# **Power System Analysis**

#### **Chapter 12 Optimal power flow**

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# Outline

- 1. Bus injection model
- 2. Branch flow model
- 3. OPF applications
- 4. Optimization algorithms

# Outline

- 1. Bus injection model
  - Single-phase devices
  - Single-phase OPF
  - OPF as QCQP
  - Three-phase devices
  - Three-phase OPF
  - Three-phase OPF as QCQP
- 2. Branch flow model
- 3. OPF applications
- 4. Optimization algorithms

# **Single-phase devices**

Voltage source *j* 

- *Ideal* voltage source: terminal voltage  $V_i$  = internal voltage
- $V_i$  is variable if the source is controllable, or given otherwise

Current source *j* 

- *Ideal* current source: terminal voltage  $I_i$  = internal voltage
- $I_i$  is variable if the source is controllable, or given otherwise

Power source j

- *Ideal* power source: terminal power  $s_j$  = internal power
- $s_j$  is variable if the source is controllable, or given otherwise

Impedance j

• Impedance  $z_j$ : constrains its terminal voltage & current  $V_j = -z_j I_j$ 

## Single-phase OPF Assumptions

Assume WLOG

- Single-phase devices: voltage sources and power sources only
- Each bus has a single device with  $\left(V_{j}, s_{j}\right)$

Formulate the simplest OPF to study general computational properties

Optimization variable:  $(V, s) := (V_j, s_j, j \in \overline{N})$ 

• Represents voltage sources  $V_j$  and power sources  $s_j$  only

Cost function  $C_0(V, s)$ 

• Fuel cost : 
$$C_0(V,s) := \sum_{j:\text{gens}} c_j \operatorname{Re}(s_j)$$

• Total real power loss: 
$$C_0(V, s) := \sum_i \operatorname{Re}(s_j)$$

Power flow equations in BIM

• Equality constraints on (*V*, *s*)

$$s_{j} = \sum_{k:j\sim k} S_{jk}(V) := \sum_{k:j\sim k} \left( y_{jk}^{s} \right)^{\mathsf{H}} \left( |V_{j}|^{2} - V_{j}V_{k}^{\mathsf{H}} \right) + \left( y_{jj}^{m} \right)^{\mathsf{H}} |V_{j}|^{2}, \qquad j \in \overline{N}$$

• Derivation:

$$I_{jk}(V) := y_{jk}^{s}(V_{j} - V_{k}) + y_{jk}^{m}V_{j}$$
  
$$S_{jk}(V) := V_{j}I_{jk}^{H}(V) := \left(y_{jk}^{s}\right)^{H}\left(|V_{j}|^{2} - V_{j}V_{k}^{H}\right) + \left(y_{jk}^{m}\right)^{H}|V_{j}|^{2}$$

- Can also use polar form and Cartesian form
- Nonlinear and global equality constraints, resulting in nonconvexity of OPF

**Operational constraints** 

- Injection limits (e.g. gen. or load capacity limits):  $s_j^{\min} \leq s_j \leq s_j^{\max}$
- Voltage limits:  $v_j^{\min} \leq |V_j|^2 \leq v_j^{\max}$
- Line limits:  $|I_{jk}(V)|^2 \le I_{jk}^{\max}$ ,  $|I_{kj}(V)|^2 \le I_{kj}^{\max}$

$$\left| \begin{array}{l} y_{jk}^{s}(V_{j} - V_{k}) + y_{jk}^{m}V_{j} \right|^{2} \leq I_{jk}^{\max}, \quad (j,k) \in E \\ \left| \begin{array}{l} y_{kj}^{s}(V_{k} - V_{j}) + y_{kj}^{m}V_{k} \right|^{2} \leq I_{kj}^{\max}, \quad (j,k) \in E \end{array} \right|$$

Line limits can also be on line powers  $(S_{jk}(V), S_{kj}(V))$  or apparent powers  $(|S_{jk}(V)|, |S_{kj}(V)|)$ 

OPF in BIM

 $\begin{array}{ll} \min_{(V,s)} & C_0(V,s) \\ \mbox{subject to} & f(V,s) = 0 & \mbox{power flow equations} \\ & g(V,s) \leq 0 & \mbox{operational constraints} \end{array}$ 

- Does not need assumption  $y_{jk}^s = y_{kj}^s$
- Can accommodate single-phase transformers with complex turns ratios

# **Single-phase OPF**

- 1. Other devices
  - Can include other devices such as current sources, impedances, capacity taps
  - Allow multiple devices connected to same bus
- 2. Can formulate OPF in terms of V only
  - Use power flow equations to express injections  $s_j(V)$  as functions of V
  - Eliminate  $s_j$  and power flow equations (equality constraints)

Next: explain each in turn

## Single-phase OPF Including other devices

Examples

- Current source (controllable): variable  $I_j$  with local constraints  $|I_j|^2 \le I_j^{\text{max}}$ ,  $s_j = V_j I_j^{\text{H}}$
- Impedance  $z_j$ : imposes additional constraint  $s_j = |V_j|^2 / z_j^H$
- Capacitor tap (controllable): variable  $y_j$  with local constraints  $y_j^{\min} \le y_j \le y_j^{\max}$ ,  $s_j = y_j^{\mathsf{H}} |V_j|^2$ • Multiple devices: injection variables  $s_{jk}$  with local constraints  $s_{jk}^{\min} \le s_{jk} \le s_{jk}^{\max}$ ,  $s_j = \sum_{jk} s_{jk}$

Including other devices at bus j imposes additional local constraints

- Additional optimization var  $u_i$  may be introduced
- Equality constraints relating  $(V_j, s_j)$  and  $u_j$  (if present) :  $f_j(V_j, s_j, u_j) = 0$
- Inequality (operational) constraints (e.g., capacity limits):  $g_i(u_i) \le 0$

#### Single-phase OPF In terms of V only

Equality constraints (BIM in complex form)

$$s_{j}(V) = \sum_{k:j \sim k} S_{jk}(V) := \sum_{k:j \sim k} \left( y_{jk}^{s} \right)^{\mathsf{H}} \left( |V_{j}|^{2} - V_{j}V_{k}^{\mathsf{H}} \right) + \left( y_{jj}^{m} \right)^{\mathsf{H}} |V_{j}|^{2}, \qquad j \in \overline{N}$$

• Expresses  $s_j$  in terms of voltages V

 $\operatorname{Cost} C_0(V) := C_0(V, s(V)) \text{ expressed as function of } V$ 

• Fuel cost:

$$C_0(V) := \sum_{j:\text{gens}} c_j \operatorname{Re}(s_j(V)) = \sum_{j:\text{gens}} c_j \operatorname{Re}\left(\sum_{k:j\sim k} \left(y_{jk}^s\right)^{\mathsf{H}} \left(|V_j|^2 - V_j V_k^{\mathsf{H}}\right) + \left(y_{jj}^m\right)^{\mathsf{H}} |V_j|^2\right)$$

• Total real power loss:

$$C_0(V) := \sum_j \operatorname{Re}(s_j(V))$$

### Single-phase OPF Operational constraints

Injection limits (e.g. generation or load capacity limits)  $s_j^{\min} \leq s_j(V) \leq s_j^{\max}$ :

$$\underline{s}_{j} \leq \sum_{k:j\sim k} \left( y_{jk}^{s} \right)^{\mathsf{H}} \left( |V_{j}|^{2} - V_{j}V_{k}^{\mathsf{H}} \right) + \left( y_{jj}^{m} \right)^{\mathsf{H}} |V_{j}|^{2} \leq \overline{s}_{j}, \qquad j \in \overline{N}$$

• Polar form:

$$\underline{p}_{j} \leq \left(\sum_{k=0}^{N} g_{jk}\right) |V_{j}|^{2} - \sum_{k \neq j} |V_{j}| |V_{k}| \left(g_{jk} \cos \theta_{jk} - b_{jk} \sin \theta_{jk}\right) \leq \overline{p}_{j}$$
$$\underline{p}_{j} \leq \left(\sum_{k=0}^{N} b_{jk}\right) |V_{j}|^{2} - \sum_{k \neq j} |V_{j}| |V_{k}| \left(b_{jk} \cos \theta_{jk} + g_{jk} \sin \theta_{jk}\right) \leq \overline{q}_{j}$$

## Single-phase OPF Operational constraints

Voltage limits (same as before):

$$v_j^{\min} \leq |V_j|^2 \leq v_j^{\max}, \quad j \in \overline{N}$$

Line limits (same as before):

$$\left| y_{jk}^{s}(V_{j} - V_{k}) + y_{jk}^{m} V_{j} \right|^{2} \leq I_{jk}^{\max}, \qquad (j,k) \in E$$

$$\left| y_{kj}^{s}(V_{k} - V_{j}) + y_{kj}^{m} V_{k} \right|^{2} \leq I_{kj}^{\max}, \qquad (j,k) \in E$$

• Line limits can also be on line powers  $\left(S_{jk}(V), S_{kj}(V)\right)$  or apparent powers  $\left(\left|S_{jk}(V)\right|, \left|S_{kj}(V)\right|\right)$ 

## Single-phase OPF In terms of V only

Feasible set

 $\mathbb{V} := \left\{ V \in \mathbb{C}^{N+1} \mid V \text{ satisfies operational constraints} \right\}$ 

OPF in BIM

 $\min_{V \in \mathbb{V}} \quad C_0(V)$ 

- Does not need assumption  $y_{jk}^s = y_{kj}^s$
- Can accommodate single-phase transformers with complex turns ratios

## Single-phase OPF In terms of V only

Feasible set

 $\mathbb{V} := \left\{ V \in \mathbb{C}^{N+1} \mid V \text{ satisfies operational constraints} \right\}$ 

OPF in BIM

 $\min_{V \in \mathbb{V}} \quad C_0(V)$ 

We will mostly study this simple OPF Can express it as a QCQP

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Quadratically constrained quadratic program:

 $\min_{x \in \mathbb{C}^n} \quad x^{\mathsf{H}} C_0 x$ 

- s.t.  $x^{\mathsf{H}}C_{l}x \leq b_{l}, \qquad l = 1, ..., L$
- $C_l: n \times n$  Hermitian matrix
- $b_l \in \mathbb{R}$
- Homogeneous QCQP : all monomials are of degree 2

Inhomogeneous QCQP

$$\min_{x \in \mathbb{C}^n} \quad x^{\mathsf{H}} C_0 x + \left( c_0^{\mathsf{H}} x + x^{\mathsf{H}} c_0 \right)$$
  
s.t. 
$$x^{\mathsf{H}} C_l x + \left( c_l^{\mathsf{H}} x + x^{\mathsf{H}} c_l \right) \leq b_l, \qquad l = 1, \dots, L$$

Homogenization:

• Idea: 
$$|x|^2 + (c^H x + x^H c) \le b \iff |x + ct|^2 - |c|^2 |t|^2 \le b, |t|^2 = 1$$

• If  $(x, t = e^{i\theta})$  satisfies 2nd inequality, then  $xt = xe^{i\theta}$  satisfies 1st inequality

Equivalent homogeneous QCQP

$$\min_{x \in \mathbb{C}^{n}, t \in \mathbb{C}} \left[ x^{\mathsf{H}} t^{\mathsf{H}} \right] \left[ \begin{matrix} C_{0} & c_{0} \\ c_{0}^{\mathsf{H}} & 0 \end{matrix} \right] \left[ \begin{matrix} x \\ t \end{matrix} \right]$$
s.t. 
$$\left[ x^{\mathsf{H}} t^{\mathsf{H}} \right] \left[ \begin{matrix} C_{l} & c_{l} \\ c_{l}^{\mathsf{H}} & 0 \end{matrix} \right] \left[ \begin{matrix} x \\ t \end{matrix} \right] \le b_{l}, \qquad l = 1, \dots, L$$

$$\left[ x^{\mathsf{H}} t^{\mathsf{H}} \right] \left[ \begin{matrix} 0 & 0 \\ 0 & 1 \end{matrix} \right] \left[ \begin{matrix} x \\ t \end{matrix} \right] = 1$$

Homogenization:

• Idea: 
$$|x|^2 + (c^H x + x^H c) \le b \iff |x + ct|^2 - |c|^2 |t|^2 \le b, |t|^2 = 1$$

• If 
$$(x, t = e^{i\theta})$$
 satisfies 2nd inequality, then  $xt = xe^{i\theta}$  satisfies 1st inequality

To write OPF as QCQP:

- Assume cost function  $C_0(V) = V^{\mathsf{H}}C_0V$  can be written as a quadratic form
- Need to rewrite operational constraints in terms of quadratic forms

**Injection limits**  $s_j^{\min} \le s_j(V) \le s_j^{\max}$ 

$$s_{j}(V) = V_{j}I_{j}^{\mathsf{H}} = \left(e_{j}^{\mathsf{H}}V\right)\left(e_{j}^{\mathsf{H}}I\right)^{\mathsf{H}} = e_{j}^{\mathsf{H}}VV^{\mathsf{H}}Y^{\mathsf{H}}e_{j}$$
$$s_{j}(V) = \operatorname{tr}\left(e_{j}^{\mathsf{H}}VV^{\mathsf{H}}Y^{\mathsf{H}}e_{j}\right) = \operatorname{tr}\left(\left(Y^{\mathsf{H}}e_{j}e_{j}^{\mathsf{H}}\right)VV^{\mathsf{H}}\right) =: V^{\mathsf{H}}Y_{j}^{\mathsf{H}}V$$

**Injection limits**  $s_j^{\min} \le s_j(V) \le s_j^{\max}$ 

$$s_{j}(V) = V_{j}I_{j}^{\mathsf{H}} = \left(e_{j}^{\mathsf{H}}V\right)\left(e_{j}^{\mathsf{H}}I\right)^{\mathsf{H}} = e_{j}^{\mathsf{H}}VV^{\mathsf{H}}Y^{\mathsf{H}}e_{j}$$
$$s_{j}(V) = \operatorname{tr}\left(e_{j}^{\mathsf{H}}VV^{\mathsf{H}}Y^{\mathsf{H}}e_{j}\right) = \operatorname{tr}\left(\left(Y^{\mathsf{H}}e_{j}e_{j}^{\mathsf{H}}\right)VV^{\mathsf{H}}\right) =: V^{\mathsf{H}}Y_{j}^{\mathsf{H}}V$$

- $Y_j$  is not Hermitian so  $V^{\mathsf{H}}Y_j^{\mathsf{H}}V$  is generally complex
- Define  $\Phi_j := \frac{1}{2} \left( Y_j^{\mathsf{H}} + Y_j \right), \qquad \Psi_j := \frac{1}{2i} \left( Y_j^{\mathsf{H}} Y_j \right)$

• Then 
$$\operatorname{Re}(s_j) = V^{\mathsf{H}} \Phi_j V$$
,  $\operatorname{Im}(s_j) = V^{\mathsf{H}} \Psi_j V$ 

Hence 
$$s_j^{\min} \le s_j(V) \le s_j^{\max}$$
 is equivalent to:  
 $p_j^{\min} \le V^{\mathsf{H}} \Phi_j V \le p_j^{\max}, \quad q_j^{\min} \le V^{\mathsf{H}} \Psi_j V \le q_j^{\max}$ 

#### **OPF as QCQP** Voltage limits

Voltage magnitude is:  $|V_j|^2 = V^H J_j V$  where  $J_j := e_j e_j^T$ 

Hence voltage limits are:  $v_j^{\min} \leq V^{\mathsf{H}} J_j V \leq v_j^{\max}$ 

## OPF as QCQP Line limits

Write  $I_{jk}$  in terms of voltage vector V:

$$I_{jk} = y_{jk}^{s}(V_{j} - V_{k}) + y_{jk}^{m}V_{j} = \left(y_{jk}^{s}(e_{j} - e_{k})^{\mathsf{T}} + y_{jk}^{m}e_{j}^{\mathsf{T}}\right)V$$

Hence current limit is:  $|I_{jk}|^2 = V^{\mathsf{H}} \hat{Y}_{jk} V \leq I_{jk}^{\max}$  where

$$\hat{Y}_{jk} := \left( y_{jk}^{s} (e_{j} - e_{k})^{\mathsf{T}} + y_{jk}^{m} e_{j}^{\mathsf{T}} \right)^{\mathsf{H}} \left( y_{jk}^{s} (e_{j} - e_{k})^{\mathsf{T}} + y_{jk}^{m} e_{j}^{\mathsf{T}} \right)$$

#### **OPF as QCQP** Simplest formulation

$$\begin{split} \min_{V \in \mathbb{C}^{N+1}} & V^{\mathsf{H}} C_0 V \\ \text{s.t.} & p_j^{\min} \leq V^{\mathsf{H}} \Phi_j V \leq p_j^{\max}, \qquad j \in \overline{N} \\ & q_j^{\min} \leq V^{\mathsf{H}} \Psi_j V \leq q_j^{\max}, \qquad j \in \overline{N} \\ & v_j^{\min} \leq V^{\mathsf{H}} J_j V \leq v_j^{\max}, \qquad j \in \overline{N} \\ & V^{\mathsf{H}} \hat{Y}_{jk} V \leq \bar{I}_{jk}^{\max}, \qquad (j,k) \in E \\ & V^{\mathsf{H}} \hat{Y}_{kj} V \leq \bar{I}_{kj}^{\max}, \qquad (j,k) \in E \end{split}$$

# Outline

#### 1. Bus injection model

- Single-phase devices
- Single-phase OPF
- OPF as QCQP
- Three-phase devices
- Three-phase OPF
- Three-phase OPF as QCQP
- 2. Branch flow model
- 3. Optimization algorithms

## Recall: overall 3-phase BIM Device + network

- 1. Device model for each 3-phase device
  - Internal model on  $\left(V_{j}^{Y/\Delta}, I_{j}^{Y/\Delta}, s_{j}^{Y/\Delta}\right)$  + conversion rules
  - External model on  $\left(V_j, I_j, s_j\right)$
  - Either can be used
  - Power source models are nonlinear; other devices are linear

#### Our perspective:

- Internal vars  $\left(V_{j}^{Y/\Delta}, I_{j}^{Y/\Delta}, s_{j}^{Y/\Delta}\right)$  are controllable, depending on types of device
- External vars  $(V_j, I_j, s_j)$  are not directly controllable
- :. use internal model + conversion rules

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## Recall: overall 3-phase BIM Device + network

- 2. Network model relates terminal vars (V, I, s)
  - Nodal current balance (linear): I = YV
  - Nodal power balance (nonlinear):

$$s_j = \sum_{k:j \sim k} \operatorname{diag} \left( V_j (V_j - V_k)^{\mathsf{H}} y_{jk}^{s\mathsf{H}} + V_j V_j^{\mathsf{H}} y_{jk}^{m\mathsf{H}} \right)$$

• Either can be used

#### For OPF, our formulation uses (V, s):

- Relate (*V*, *s*) through power flow equations
- Power sources lead to nonlinear analysis, even if we use I = YV as network equation
- Need to relate internal optimization vars to  $(V_j, s_j)$  using conversion rules

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# **Three-phase devices**

Voltage source  $V_i^{Y/\Delta}$ 

- Internal optimization variable  $u_j := V_j^{Y/\Delta}$  ( $\gamma_j^Y$  assumed given)
- Local constraints that relate internal vars to  $(V_j, s_j)$

$$Y: V_j = V_j^Y + \gamma_j^Y \mathbf{1}$$
  
$$\Delta: \Gamma V_j = V_j^{\Delta}$$

#### Note:

- Choosing  $V_j^{\Delta}$  does not uniquely determine  $V_j$
- Optimization over  $V_j$  implicitly chooses an optimal  $\gamma_j^{\Delta} := \frac{1}{3} \mathbf{1}^{\mathsf{T}} V_j$
- If  $\gamma_j^\Delta$  is given, then  $\Gamma V_j = V_j^\Delta$  should be replaced by  $V_j = \Gamma^\dagger V_j^\Delta + \gamma_j^\Delta$ 1

# **Three-phase devices**

Current source  $I_i^{Y/\Delta}$ 

- Internal optimization variable  $u_j := I_j^{Y/\Delta}$
- Local constraints that relate internal vars and  $\left(V_{j}, s_{j}\right)$

$$\begin{array}{lll}Y: & s_{j} &= -\operatorname{diag}\left(V_{j}I_{j}^{Y\mathsf{H}}\right)\\ \Delta: & s_{j} &= -\operatorname{diag}\left(V_{j}I_{j}^{\Delta\mathsf{H}}\Gamma\right)\end{array}$$

#### Note:

- Optimization over  $I_j^{\Delta}$  implicitly chooses an optimal  $\beta_j^{\Delta} := \frac{1}{3} \mathbf{1}^T I_j^{\Delta}$
- If  $\beta_j^{\Delta}$  is given, it imposes an additional constraint  $I_j^{\Delta} = -\frac{1}{3}\Gamma I_j + \beta_j^{\Delta} 1$  (and express  $I_j$  in terms of  $(V_j, s_j)$

## **Three-phase devices**

Power source  $\left(s_{j}^{Y/\Delta}, I_{j}^{Y/\Delta}\right)$ 

• Internal optimization variable  $u_j := \left(s_j^{Y/\Delta}, I_j^{Y/\Delta}\right)$  (assume  $\gamma_j^Y := V_j^n = 0$ )

• Local constraints that relate internal vars and  $(V_j, s_j)$ 

$$\begin{split} Y: & s_j = -s_j^Y \\ \Delta: & s_j = -\operatorname{diag}\left(V_j I_j^{\Delta \mathsf{H}} \Gamma\right), \qquad s_j^{\Delta} = \operatorname{diag}\left(\Gamma V_j I_j^{\Delta \mathsf{H}}\right) \end{split}$$

Impedance  $z_i^{Y/\Delta}$ 

- Given parameter:  $z_j^{Y/\Delta}$  (assume  $\gamma_j^Y := V_j^n = 0$ )
- Local constraints on terminal vars  $(V_j, s_j)$

# **Three-phase OPF**

Variables:

- Terminal variables  $\left(V_{j}, s_{j}\right)$
- Internal variables  $u_i$  depending on devices (discussed above)

Cost function:  $C_0(V, s, u)$ 

Equality constraints:

1. Power flow equations on (V, s) (global constraint): f(V, s) = 0

$$s_j = \sum_{k:j \sim k} \operatorname{diag}\left(V_j(V_j - V_k)^H \left(y_{jk}^s\right)^H + V_j V_j^H \left(y_{jk}^m\right)^H\right), \quad j \in \overline{N}$$

2. Conversion rules relating internal optimization var  $u_j$  to  $(V_j, s_j)$  (local constraint, discussed above)

$$f_j^{Y/\Delta}\left(V_j, s_j, u_j\right) = 0, \qquad j \in \overline{N}$$

# **Three-phase OPF**

#### Inequality constraints:

1. Operational constraints on external vars:  $g(V, s) \le 0$ 

injection limits:
$$s_j^{\phi \min} \leq s_j^{\phi} \leq s_j^{\phi \max}$$
, $\phi \in \{a, b, c\}$ ,  $j \in \overline{N}$ voltage limits: $v_j^{\phi \min} \leq \left| V_j^{\phi} \right|^2 \leq v_j^{\phi \max}$ , $\phi \in \{a, b, c\}$ ,  $j \in \overline{N}$ line limits: $\left| I_{jk}^{\phi}(V) \right|^2 \leq I_{jk}^{\phi \max}$ , $\left| I_{kj}^{\phi}(V) \right|^2 \leq I_{kj}^{\phi \max}$ , $\phi \in \{a, b, c\}$ ,  $(j, k) \in E$ 

Same constraints as single-phase OPF, but on single-phase equivalent circuit

# **Three-phase OPF**

Inequality constraints:

2. Operational constraints on internal vars:  $g_j^{Y/\Delta}(u_j) \le 0$ for  $\phi n \in \{an, bn, cn\}$ ,  $\phi \varphi \in \{ab, bc, ca\}$ voltage source:  $v_j^{\phi n \min} \le |V_j^{\phi n}|^2 \le v_j^{\phi n \max}$ ,  $v_j^{\phi \varphi \min} \le |V_j^{\phi \varphi}|^2 \le v_j^{\phi \varphi \max}$ current source:  $|I_j^{\phi n}|^2 \le I_j^{\max}$ ,  $|I_j^{\phi \varphi}|^2 \le I_j^{\max}$ power source:  $s_j^{\phi n \min} \le s_j^{\phi n} \le s_j^{\phi n \max}$ ,  $|I_j^{\phi \varphi}|^2 \le I_j^{\phi n \max}$  $s_j^{\phi \varphi \min} \le s_j^{\phi \varphi} \le s_j^{\phi \varphi \max}$ ,  $|I_j^{\phi \varphi}|^2 \le I_j^{\phi \varphi \max}$ 

Local constraints at each bus j

#### Three-phase OPF Constraints summary

1. Constraints on terminal variables: f(V, s) = 0,  $g(V, s) \le 0$ 

- Power flow equation and operational constraints (terminal power injection limits, voltage limits, line limits)
- Global constraints
- Extension of single-phase constraints to 3-phase setting, using single-phase equivalent
- 2. Conversion rules relating  $u_j$  and  $\left(V_j, s_j\right)$ :  $f_j^{Y/\Delta}\left(u_j, V_j, s_j\right) = 0$ 
  - Local equality constraint for each device j
- 3. Operational constraints on internal variables:  $g_j^{Y/\Delta}(u_j) \le 0$ 
  - Depending on type of device (voltage and capacity limits)
  - Local constraints for each device *j*
## Three-phase OPF Simplest formulation

#### OPF in BIM

$$\begin{array}{ll} \min_{(V,s,u)} & C_0(V,s,u) \\ f(V,s) = 0, & g(V,s) \leq 0 \\ f_j^{Y/\Delta}(V_j,s_j,u_j) = 0, & g_j^{Y/\Delta}(u_j) \leq 0, & j \in \overline{N} \end{array} \begin{array}{ll} \mbox{Global constraints on terminal vars} \\ \mbox{Local constraints at each bus } j \end{array}$$

## Three-phase OPF As QCQP

- 1. Can formulate OPF in terms of (V, u) only
  - Use power flow equations to express  $s_j(V) = V^H \left( Y_j^{\phi H} \right) V$  and eliminate  $s_j$  and f(V, s) = 0
  - Same idea as before applied to single-phase equivalent
- 2. Can formulate OPF as QCQP
  - Express operational constraints  $g(V, s(V)) \le 0$  in terms of quadratic forms in V (same idea applied to single-phase equivalent)
  - Express conversion rules  $f_j^{Y/\Delta}(V_j, s_j(V), u_j) = 0$  in terms of quadratic forms in  $(V, u_j)$

#### For details: see Lecture Notes

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## **Overview**



## **Assumptions** Both single-phase & 3-phase OPF

Radial network

• BFM most useful for modeling distribution systems

$$z_{jk}^s = z_{kj}^s$$
 or equivalently  $y_{jk}^s = y_{kj}^s$ 

• Does not include 3-phase transformers in  $\Delta Y$  or  $Y\Delta$  configuration (or single-phase transformers with complex gains)

 $y_{jk}^m = y_{kj}^m = 0$ 

• Reasonable assumption for distribution line where  $|y_{ik}^m|, |y_{kj}^m| \ll |y_{ik}^s|$ 

Includes only voltage sources and power sources

- Optimization variables are voltages (squared magnitudes)  $v_i$  and power injections  $s_i$  respectively
- A current source or an impedance will introduce additional var and constraint.

## Single-phase OPF

Power flow equations

• All lines point away from bus 0 (root)

$$\sum_{k:j \to k} S_{jk} = S_{ij} - z_{ij} \ell_{ij} + s_j, \quad j \in \overline{N}$$
$$v_j - v_k = 2 \operatorname{Re} \left( z_{jk}^{\mathsf{H}} S_{jk} \right) - |z_{jk}|^2 \ell_{jk}, \quad j \to k \in E$$
$$v_j \ell_{jk} = |S_{jk}|^2, \quad j \to k \in E$$

**Operational constraints** 

$$s_{j}^{\min} \leq s_{j} \leq s_{j}^{\max}$$
$$v_{j}^{\min} \leq v_{j} \leq v_{j}^{\max}$$
$$\ell_{jk} \leq I_{jk}^{\max}$$

Steven Low OPF Branch flow model

# Single-phase OPF

Feasible set

 $\mathbb{T}_0 := \left\{ x := (s, v, \ell, S) \in \mathbb{R}^{6N+3} \mid x \text{ satisfies PF equations & operational constraints} \right\}$ 

OPF in BFM

 $\min_{x \in \mathbb{T}_0} \quad C(x)$ 

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## Single-phase OPF Equivalence

Recall for BIM:

- Feasible set:  $\mathbb{V} := \{ V \in \mathbb{C}^{N+1} \mid V \text{ satisfies operational constraints} \}$
- $\bullet \ \ {\rm OPF:} \ \min_{V \in \mathbb{V}} \quad C_0(V)$

OPF in BFM is equivalent to OPF in BIM:

- Feasible sets  $\mathbb{T}_0$  and  $\mathbb V$  are equivalent (Ch 6)
- ... provided cost functions C(x) and  $C_0(V)$  are the same

Variables (x, u):

1. Directly generalizes vars in single-phase OPF ( $\mathbb{S}^n_+$ : complex psd matrices)

$$s_{j} \in \mathbb{C}^{3}, \qquad v_{j} \in \mathbb{S}^{3}_{+}, \qquad j \in \overline{N}$$
$$\ell_{jk} \in \mathbb{S}^{3}_{+}, \qquad S_{jk} \in \mathbb{C}^{3 \times 3}, \qquad j \to k \in E$$

To write conversion rule for power sources, introduce phasors as additional vars

$$\left(V_{j}, j \in \overline{N}\right), \qquad \left(\tilde{I}_{jk}, j \to k \in E\right)$$

Let  $x := (s, v, \ell, V, \tilde{I}, S)$ 

Variables (x, u):

2. Internal variables  $u := (u_j, j \in \overline{N})$  of 3-phase devices

voltage source :

power source :

Equality constraints

1. Power flow equations (from Ch 10):

$$\begin{split} \sum_{k:j \to k} \operatorname{diag}(S_{jk}) &= \operatorname{diag}\left(S_{ij} - z_{ij}\ell_{ij}\right) + s_j, \qquad j \in \overline{N} \\ v_j - v_k &= \left(z_{jk} S_{jk}^{\mathsf{H}} + S_{jk} z_{jk}^{\mathsf{H}}\right) - z_{jk} \ell_{jk} z_{jk}^{\mathsf{H}}, \qquad j \to k \in E \\ \begin{bmatrix} v_j S_{jk} \\ S_{jk}^{\mathsf{H}} \ell_{jk} \end{bmatrix} \geq 0, \qquad j \to k \in E \\ \operatorname{rank} \begin{bmatrix} v_j S_{jk} \\ S_{jk}^{\mathsf{H}} \ell_{jk} \end{bmatrix} = 1, \qquad j \to k \in E \\ v_j &= V_j V_j^{\mathsf{H}}, \quad \ell_{jk} = \tilde{I}_{jk} \tilde{I}_{jk}^{\mathsf{H}}, \quad S_{jk} = V_j \tilde{I}_{jk}^{\mathsf{H}}, \qquad j \to k \in E \end{split}$$
 additional equations

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Equality constraints

1. Power flow equations (from Ch 10):

$$\begin{split} \sum_{k:j \to k} \operatorname{diag}(S_{jk}) &= \operatorname{diag}\left(S_{ij} - z_{ij} \ell_{ij}\right) + s_j, \quad j \in \overline{N} \\ v_j - v_k &= \left(z_{jk} S_{jk}^{\mathsf{H}} + S_{jk} z_{jk}^{\mathsf{H}}\right) - z_{jk} \ell_{jk} z_{jk}^{\mathsf{H}}, \quad j \to k \in E \\ \begin{bmatrix} v_j S_{jk} \\ S_{jk}^{\mathsf{H}} \ell_{jk} \end{bmatrix} \geq 0, \quad j \to k \in E \\ \operatorname{redundant\ constraints\ kept\ for\ semidefinite\ relaxation\ (later)} \\ \operatorname{rank} \begin{bmatrix} v_j S_{jk} \\ S_{jk}^{\mathsf{H}} \ell_{jk} \end{bmatrix} = 1, \quad j \to k \in E \\ v_j &= V_j V_j^{\mathsf{H}}, \quad \ell_{jk} = \tilde{I}_{jk} \tilde{I}_{jk}^{\mathsf{H}}, \quad S_{jk} = V_j \tilde{I}_{jk}^{\mathsf{H}}, \quad j \to k \in E \end{split}$$

Equality constraints

2. Conversion rules for voltage & power sources (assume  $\gamma_j^Y := V_j^n = 0$ )

voltage source : 
$$Y$$
:  $v_j = V_j^Y V_j^{YH} = u_j u_j^H$   
 $\Delta$ :  $\Gamma v_j \Gamma^T = V_j^\Delta V_j^{\Delta H} = u_j u_j^H$ 

power source : 
$$Y$$
:  $s_j = -\operatorname{diag}\left(V_j u_{j2}^{\mathsf{H}}\right), \qquad s_j = -u_{j1}$   
 $\Delta$ :  $s_j = -\operatorname{diag}\left(V_j u_{j2}^{\mathsf{H}}\Gamma\right), \qquad u_{j1} = \operatorname{diag}\left(\Gamma V_j u_{j2}^{\mathsf{H}}\right)$ 

#### Inequality constraints

1. Operational constraints on *x*:

injection limits: $s_j^{\min} \leq s_j \leq s_j^{\max}$ , $j \in \overline{N}$ voltage limits: $v_j^{\min} \leq \text{diag}\left(v_j\right) \leq v_j^{\max}$ , $j \in \overline{N}$ line limits: $\text{diag}\left(\ell_{jk}\right) \leq I_{jk}^{\max}$ , $(j,k) \in E$ 

### Inequality constraints

2. Operational constraints on internal vars  $u_j$ :

$$\begin{array}{ll} \text{voltage source:} & v_j^{\phi n \min} \leq \left| V_j^{\phi n} \right|^2 \leq v_j^{\phi n \max}, & v_j^{\phi \varphi \min} \leq \left| V_j^{\phi \varphi} \right|^2 \leq v_j^{\phi \varphi \max} \\ \text{power source:} & s_j^{Y \min} \leq s_j^Y \leq s_j^{Y \max}, & \left| I_j^{\phi n} \right|^2 \leq I_j^{\phi n \max} \\ & s_j^{\Delta \min} \leq s_j^{\Delta} \leq s_j^{\Delta \max}, & \left| I_j^{\phi \varphi} \right|^2 \leq I_j^{\phi \varphi \max} \end{array}$$

Feasible set

$$\mathbb{T}_{3p} := \left\{ (x, u) := (s, v, \ell, V, \tilde{I}, S, u) \mid (x, u) \text{ satisfies all constraints} \right\}$$

OPF in BFM

 $\min_{(x,u)\in\mathbb{T}_{3p}} \quad C(x,u)$ 

Three-phase OPF in BFM is equivalent to three-phase OPF in BIM:

- Their feasible sets are equivalent (Ch 10)
- ... provided their cost functions are equivalent

# Outline

- 1. Bus injection model
- 2. Branch flow model
- 3. OPF applications
  - Voltage control (distribution grid)
- 4. Optimization algorithms

## Voltage control Distribution system

Voltage instability: magnitudes fluctuate outside their limits

- PVs may push magnitudes above upper limits
- EVs may push magnitudes below lower limits

Traditional solution

• Infrastructure upgrade: more/larger transformers, wires, etc

Non-wire solution

- Distributed energy resources (DER) optimization
- e.g. batteries, smart inverters, demand response
- Can formulate as an OPF

**Optimal battery operation** 

$$\min_{u,V, b} \sum_{t} \sum_{j} \left( |V_{j}(t)|^{2} - v_{j}^{\mathsf{ref}}(t) \right)^{2}$$
 deviation from nominal voltages  
s.t.  $u_{j}(t) + \sigma_{j}(t) = \sum_{k:j \sim k} S_{jk}(V(t)), \quad \underline{v}_{j} \leq |V_{j}(t)|^{2} \leq \overline{v}_{j}$   
 $|S_{jk}(V(t))| \leq \overline{S}_{jk}, \quad |S_{kj}(V(t))| \leq \overline{S}_{kj}$ 

**Optimal battery operation** 

$$\min_{u,V, b} \sum_{t} \sum_{j} \left( |V_{j}(t)|^{2} - v_{j}^{\text{ref}}(t) \right)^{2}$$
 deviation from nominal voltages  
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 $|S_{jk}(V(t))| \leq \overline{S}_{jk}, \quad |S_{kj}(V(t))| \leq \overline{S}_{kj}$   
 $b_{j}(t+1) = b_{j}(t) - \operatorname{Re}\left(u_{j}(t)\right)$  charging/discharging (100% efficiency)

**Optimal battery operation** 

$$\min_{u,V, b} \sum_{t} \sum_{j} \left( |V_{j}(t)|^{2} - v_{j}^{\text{ref}}(t) \right)^{2}$$
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 $|S_{jk}(V(t))| \leq \overline{S}_{jk}, \quad |S_{kj}(V(t))| \leq \overline{S}_{kj}$   
 $b_{j}(t+1) = b_{j}(t) - \operatorname{Re}\left(u_{j}(t)\right)$  charging/discharging (100% efficiency)  
 $\underline{u}_{j} \leq \operatorname{Re}\left(u_{j}(t)\right) \leq \overline{u}_{j}, \quad 0 \leq b_{j}(t) \leq B_{j}$   
power limit energy limit

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**Optimal battery placement** 

$$\begin{split} \min_{u,V, b, B} & \sum_{t} \sum_{j} \left( |V_{j}(t)|^{2} - v_{j}^{\mathsf{ref}}(t) \right)^{2} + \sum_{j} c_{j}B_{j} \\ \text{s.t.} & u_{j}(t) + \sigma_{j}(t) = \sum_{k:j \sim k} S_{jk}(V(t)), \quad \underline{v}_{j} \leq |V_{j}(t)|^{2} \leq \overline{v}_{j} \\ & |S_{jk}(V(t))| \leq \overline{S}_{jk}, \quad |S_{kj}(V(t))| \leq \overline{S}_{kj} \\ & b_{j}(t+1) = b_{j}(t) - \operatorname{Re}\left(u_{j}(t)\right) \\ & \underline{u}_{j} \leq \operatorname{Re}\left(u_{j}(t)\right) \leq \overline{u}_{j}, \quad 0 \leq b_{j}(t) \leq B_{j} \\ & B_{j}^{\mathsf{opt}} > 0 : \mathsf{place battery at bus } j \end{split}$$

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# Outline

- 1. Bus injection model
- 2. Branch flow model
- 3. Optimization algorithms
  - Newton-Raphson algorithm
  - Interior-point algorithm

# **Complex formulation**

Even though OPF is often formulated in  $\mathbb{C}$ , it is converted to  $\mathbb{R}$  before being solved iteratively

### **Example: QCQP**

- $\min_{x \in \mathbb{C}^n} x^{\mathsf{H}} C_0 x$ s.t.  $x^{\mathsf{H}} C_l x \leq b_l, \qquad l = 1, \dots, L$
- $C_l: n \times n$  Hermitian matrix
- $b_l \in \mathbb{R}$

#### **Equivalent to:**

$$\min_{\substack{(x_r,x_i)\in\mathbb{R}^{2n}\\ \text{s.t.}}} \begin{bmatrix} x_r\\x_i \end{bmatrix}^{\mathsf{T}} \begin{bmatrix} C_{0r} & -C_{0i}\\ C_{0i} & C_{0r} \end{bmatrix} \begin{bmatrix} x_r\\x_i \end{bmatrix}$$
$$s.t. \begin{bmatrix} x_r\\x_i \end{bmatrix}^{\mathsf{T}} \begin{bmatrix} C_{lr} & -C_{li}\\ C_{li} & C_{lr} \end{bmatrix} \begin{bmatrix} x_r\\x_i \end{bmatrix} \le b_l, \quad l = 1, \dots, L$$

•  $2n \times 2n$  symmetric matrices

## Algorithms for OPF Popular algorithms

Newton-Raphson algorithm

- 2nd order algorithm
- Interior-point algorithm

Interior-point algorithm

- Based on barrier functions
- Uses of Newton-Raphson algorithm for subproblems

# **Newton-Raphson algorithm**

NR is algorithm for solving

$$F(x) = 0, \qquad F: \mathbb{R}^n \to R^n$$

Iteratively:

$$x(t+1) = y(t) + \Delta x(t)$$
$$J(y(t)) \Delta x(t) = -F(x(t))$$
where  $J(x) := \frac{\partial F}{\partial x}(x)$  is Facobian of  $F$ 

Application to optimization problems:

- F(x) = 0 is KKT condition
- If NR converges, it computes a KKT point  $x^{opt}$
- *x*<sup>opt</sup> is a global optimal if the problem is convex (feasible otherwise)

# **Newton-Raphson algorithm**

Describe NR progressively for solving

- Linear equality constrained problems
- Nonlinear equality constrained problems
- Inequality constrained problems

## Newton-Raphson algorithm Linear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \quad \text{s.t.} \quad Ax = b$ 

where

- $f: \mathbb{R}^n \to \mathbb{R}$  is twice continuously differentiable
- $A \in \mathbb{R}^{m \times n}$

### Newton-Raphson algorithm Linear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \quad \text{s.t.} \quad Ax = b$ 

Lagrangian:

$$L(x,\lambda) := f(x) + \lambda^{\mathsf{T}}(Ax - b)$$

Jacobian of  $L(x, \lambda)$ :

$$F(x,\lambda) := \begin{bmatrix} \nabla_x L(x,\lambda) \\ \nabla_\lambda L(x,\lambda) \end{bmatrix} = \begin{bmatrix} \nabla f(x) + A^{\mathsf{T}}\lambda \\ Ax - b \end{bmatrix}$$

KKT condition to be solved by NR algorithm:

 $F(x,\lambda)=0$ 

## Newton-Raphson algorithm Linear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \quad \text{s.t.} \quad Ax = b$ 

Jacobian of  $F(x, \lambda)$ :

$$J(x,\lambda) = \begin{bmatrix} \frac{\partial^2 f}{\partial x^2}(x) & A^{\mathsf{T}} \\ A & 0 \end{bmatrix}$$
 • KKT matrix  
• Independent of  $\lambda$ 

NR iteration:

$$\begin{bmatrix} x(t+1)\\\lambda(t+1) \end{bmatrix} = \begin{bmatrix} x(t)\\\lambda(t) \end{bmatrix} + \begin{bmatrix} \Delta x(t)\\\Delta\lambda(t) \end{bmatrix} \text{ where } \begin{bmatrix} \frac{\partial^2 f}{\partial x^2}(x(t)) & A^{\mathsf{T}}\\A & 0 \end{bmatrix} \begin{bmatrix} \Delta x(t)\\\Delta\lambda(t) \end{bmatrix} = -\begin{bmatrix} \nabla f(x(t)) + A^{\mathsf{T}}\lambda(t)\\Ax(t) - b \end{bmatrix}$$

## Newton-Raphson algorithm Nonlinear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \qquad \text{s.t.} \qquad g(x) = 0$ 

where

•  $f: \mathbb{R}^n \to \mathbb{R}$  and  $g: \mathbb{R}^n \to \mathbb{R}^m$  are twice continuously differentiable

Follow the same procedure as for linear equality constrained problems

### Newton-Raphson algorithm Nonlinear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \qquad \text{s.t.} \qquad g(x) = 0$ 

Lagrangian:

$$L(x,\lambda) := f(x) + \lambda^{\mathsf{T}} g(x)$$

Jacobian of  $L(x, \lambda)$ :

$$F(x,\lambda) := \begin{bmatrix} \nabla_x L(x,\lambda) \\ \nabla_\lambda L(x,\lambda) \end{bmatrix} = \begin{bmatrix} \nabla f(x) + \frac{\partial g}{\partial x}(x)^{\mathsf{T}} \lambda \\ g(x) \end{bmatrix}$$

KKT condition to be solved by NR algorithm:

 $F(x,\lambda) = 0$ 

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### Newton-Raphson algorithm Nonlinear equality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \qquad \text{s.t.} \qquad g(x) = 0$ 

Jacobian of  $F(x, \lambda)$ :

$$J(x,\lambda) = \begin{bmatrix} \frac{\partial^2 f}{\partial x^2}(x) + \sum_k \frac{\partial^2 g_k}{\partial x^2} \lambda_k & \frac{\partial g}{\partial x}(x)^\mathsf{T} \\ \frac{\partial g}{\partial x}(x) & 0 \end{bmatrix}$$

NR iteration:

$$\begin{bmatrix} x(t+1) \\ \lambda(t+1) \end{bmatrix} = \begin{bmatrix} x(t) \\ \lambda(t) \end{bmatrix} + \begin{bmatrix} \Delta x(t) \\ \Delta \lambda(t) \end{bmatrix} \text{ where } J(x,\lambda) \begin{bmatrix} \Delta x(t) \\ \Delta \lambda(t) \end{bmatrix} = - \begin{bmatrix} \nabla f(x(t)) + \frac{\partial g}{\partial x}(x(t))^{\mathsf{T}} \lambda(t) \\ g(x(t)) \end{bmatrix}$$

## Newton-Raphson algorithm Inequality constraint

Consider

 $\min_{x \in \mathbb{R}^n} f(x) \qquad \text{s.t.} \qquad g(x) \le 0$ 

where

•  $f: \mathbb{R}^n \to \mathbb{R}$  and  $g: \mathbb{R}^n \to \mathbb{R}^m$  are twice continuously differentiable

#### Two common solution approaches

1. Introduce slack var  $z \ge 0$  to reduce the inequality into a simple inequality constraint:

 $\min_{(x,z)\in\mathbb{R}^{n+m}} f(x) \qquad \text{s.t.} \qquad g(x)+z = 0, \quad z \ge 0$ 

2. Replace constraint by a penalty term and reduce to unconstrained problem:

$$\min_{x \in \mathbb{R}^n} \quad f(x) + \frac{1}{t}\phi(x)$$

This is the approach of interior-point algorithms !

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## Interior-point algorithm Basic idea

Consider

 $\min_{x \in \mathbb{R}^n} f_0(x) \qquad \text{s.t.} \qquad f(x) \le 0, \quad g(x) = 0$ 

where

•  $f_0: \mathbb{R}^n \to \mathbb{R}, f: \mathbb{R}^n \to \mathbb{R}^m, g: \mathbb{R}^n \to \mathbb{R}^p$  are twice continuously differentiable

#### **Basic idea:**

- Approximate problem by equality constrained problem by replacing  $f(x) \le 0$  by a barrier function
- Solve the approximate problem by Newton-Raphson methods

### **Interior-point algorithm** Log barrier function

Log barrier function 
$$\phi : \mathbb{R}^n \to \mathbb{R}$$
 is  
 $\phi(x) := -\sum_{i=1}^m \log(-f_i(x))$   
over  $\operatorname{dom} \phi := \{x \in \mathbb{R}^n : f_i(x) < 0, i = 1, ..., m\}$ 

#### **Properties:**

• 
$$\phi(x) \to \infty \text{ as } f_i(x) \to 0 \text{ for any } i$$
  
•  $\nabla \phi(x) = \sum_{i=1}^m \frac{1}{-f_i(x)} \nabla f_i(x)$   
•  $\frac{\partial^2 \phi}{\partial x^2}(x) = \sum_i \frac{1}{f_i^2(x)} \nabla f_i(x) \nabla f_i^{\mathsf{T}}(x) + \sum_i \frac{1}{-f_i(x)} \frac{\partial^2 f_i}{\partial x^2}(x)$ 

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### Interior-point algorithm Approximate problem

Consider

 $\min_{x \in \mathbb{R}^n} f_0(x) \quad \text{s.t.} \quad f(x) \le 0, \quad g(x) = 0$ 

#### Approximate problem

$$\min_{x \in \mathbb{R}^n} f_0(x) + \frac{1}{t}\phi(x) \qquad \text{s.t.} \qquad g(x) = 0$$

or

Problem(t) :

$$\min_{x \in \mathbb{R}^n} tf_0(x) + \phi(x) \qquad \text{s.t.} \qquad g(x) = 0$$

• Larger  $t > 0 \implies$  more accurate approximation

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# **Barrier method**

### A popular interior-point method

#### Basic idea

- Solve Problem(*t*) for an increasing sequence of t > 0 until solution is accurate enough
- For each *t*, solve Problem(*t*) using Newton-Raphson algorithm

#### <u>Questions</u>

- How to choose the sequence of *t*?
- When to terminate?

Answer these question for convex problems

### **Barrier method** Assumptions

- 1. Original problem is convex, i.e.,  $f_0, f_1, \ldots, f_m$  are convex and g(x) = Ax b
- 2. For each t > 0, Newton-Raphson algorithm converges to the unique optimal solution x(t) of the approximate problem
- Central point : optimal solution *x*(*t*)
- Central path : set  $\{x(t) : t > 0\}$  of central points

### **Barrier method** Central point *x*(*t*)

- 1. Original problem is convex, i.e.,  $f_0, f_1, \dots, f_m$  are convex and g(x) = Ax b
- 2. For each t > 0, Newton-Raphson algorithm converges to the unique optimal solution x(t) of the approximate problem

#### **Theorem**

For each t > 0

- 1. x(t) is feasible for original problem
- 2. Objective value is at most m/t away from optimal value, i.e.,  $f_0(x(t)) f_0^{\text{opt}} \leq \frac{m}{t}$ In particular  $f_0(x(t)) \to f_0^{\text{opt}}$  as  $t \to \infty$

## **Barrier method**

**Input:** *strictly* feasible *x*, initial  $t := t_0$ , scaling factor  $\gamma > 1$ , tolerance  $\varepsilon$ . **Output:** an approximate solution *x* 

- 1. while  $t \leq \frac{m}{\epsilon}$  do
  - (a) Solve Problem(t) to compute x(t) using the Newton-Raphson algorithm starting from x.
    (b) x ← x(t).
    (c) t ← γt.
- 2. **Return**: *x*.

In principle, one can solve Problem(t) with  $t := m/\epsilon$  instead of solving a sequence of Problem(t). In practice, barrier method works better.

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